

Subroutine ATEIG

This subroutine computes the eigenvalues of a real upper almost-triangular matrix (Hessenberg form -- see subroutine HSBG) using the double QR iteration of J. G. F. Francis.

1. Mathematical background

a. Definition of the QR iteration

Let A be a real or complex nonsingular matrix of order n . Then a decomposition of A exists of the form

$$A = QR$$

where Q is unitary and R is upper triangular. If the diagonal elements of R are real and positive, Q is unique. Consider now the sequence of matrices $A^{(p)}$ defined recursively by

$$A^{(0)} = A, A^{(p)} = Q^{(p)} R^{(p)}, A^{(p+1)} = R^{(p)} Q^{(p)} \quad p \geq 0.$$

Note that $A^{(p+1)} = Q^{(p)*} A^{(p)} Q^{(p)}$ for $p \geq 0$; hence it follows that $A^{(p)}$ is similar to A for all p .

Furthermore, if A satisfies certain conditions, it can be proved that $A^{(p)}$ tends to an upper triangular matrix as $p \rightarrow \infty$; thus the eigenvalues of A are the diagonal elements of this limit matrix.

b. Convergence

If the moduli of the eigenvalues are distinct, the elements $a_{ij}^{(p)}$ below the main diagonal of $A^{(p)}$ tend to zero, as do $|\lambda_i|^{p/|\lambda_j|}$, the eigenvalues being subscripted so that $|\lambda_i| > |\lambda_{i+1}|$.

Thus, in general, the eigenvalues appear on the main diagonal, starting from the last position, in increasing order of moduli.

So, when the smallest eigenvalue λ_n has been found, we can reduce the order of the matrix by neglecting the last row and column and find λ_{n-1} by the same process, without any special deflation.

Note that the speed of convergence is considerably improved when the origin of the eigenvalues is shifted close to λ_n .

Such a shift, say $s^{(p)}$, can be introduced before an iteration and the opposite one afterwards. Then the iteration can be written as:

$$A^{(p)} - s^{(p)} I = Q^{(p)} R^{(p)}$$

$$A^{(p+1)} = R^{(p)} Q^{(p)} + s^{(p)} I$$

In general, $A_{n,n}^{(p)}$, for p large enough, can provide an efficient value for $s^{(p)}$.

c. Use of the Hessenberg form

The Hessenberg form is preserved under the QR iteration. Thus, a reduction of the initial matrix to the Hessenberg form can give a significant saving of computation in each iteration for the QR decomposition, the lower part of the matrix consisting only of the codiagonal terms.

Before each iteration, the codiagonal terms will be inspected. If some of these are zero, the matrix will be split according to this occurrence, and the iteration will be applied to the lower main submatrix only.

d. The double QR iteration

Let A be a diagonalizable real upper Hessenberg matrix. Such a matrix must be expected to have complex conjugate pairs of eigenvalues. If these pairs are the only eigenvalues of equal modulus, it can be shown that they will appear as the latent roots of main submatrices of order 2. In this case, if a shift is close to one of these roots, it will be complex, and we will have to deal with complex matrices, although the initial one is real. The use of the double QR iteration avoids this inconvenience.

Taking $s^{(p+1)} = \bar{s}^{(p)}$, consider the transformation giving $A^{(p+2)}$ from $A^{(p)}$:

$$A^{(p+2)} = Q^{(p+1)*} Q^{(p)*} A^{(p)} Q^{(p)} Q^{(p+1)}$$

It can be proved that the product $Q^{(p)} Q^{(p+1)}$ derives from the QR decomposition of the matrix $M = (A^{(p)} - s^{(p)} I) (A^{(p)} - \bar{s}^{(p+1)} I)$, which is real.

In fact, Francis (1961, 1962) showed that only the first column m_1 of M is necessary for determining the transformation which gives $A^{(p+2)}$ from $A^{(p)}$, if they both have the Hessenberg form.

Practically, the first part of the double iteration consists of the application of an initial transformation $N_1^* A^{(p)} N_1$ where N_1 is unitary and such that $N_1^* m_1 = \pm \|m_1\| e_1$. This leads to a matrix which no longer has the Hessenberg form.

Thus, the remaining part of the iteration will involve the application of $(n-1)$ successive transformations, which have the same form as the initial one whose matrices N_i are such that the resulting matrix $A^{(p+2)}$ has the Hessenberg form.

This process can fail when a subdiagonal term of the given matrix is zero. In this case, the matrix can be split, and the iteration is performed on the lower main submatrix only.

In the subroutine, N_i are Householder's matrices.

2. Programming considerations

At each iteration, the latent roots x_1 and x_2 of the lower main submatrix of order 2 are computed. Then, the following situations can occur:

- The term $a_{n-1,n-2}$ can be taken as zero. Then, x_1 and x_2 are eigenvalues of the original matrix, and the order of the matrix is reduced by 2. IANA(N) and IANA(N-1) are set to 0 and 2 respectively.
- The term $a_{n,n-1}$ can be taken as zero. In this case, $a_{n,n}$ is an eigenvalue of the original matrix, and the order of the matrix is reduced by 1. IANA(N) is set to 1.
- One of the last two subdiagonal terms is stable through one iteration. Then the smaller one is considered as zero. The corresponding components of IANA are set to 0, 1, or 2, according to a. or b.
- The maximum number of iterations is reached. In this case, the smaller of the last two subdiagonal elements is taken as zero. The corresponding components of IANA are set to 0, 1, or 2, according to a. or b.

The user can check the results by inspecting the subdiagonal terms of the matrix on return from the subroutine, according to the vector IANA, in the following way:

If for each IANA(I) containing 1 or 2, $2 \leq I \leq M$,

$$|A(I, I-1)| \leq 10^{-7} (|RR(I)| + |RI(I)|),$$

then RR(I) and RI(I) were computed with a satisfactory accuracy.

For reference see:

- J.G.F. Francis, Computer Journal. October, 1961 4-3, January, 1962 4-4.
- J.H. Wilkinson, The Algebraic Eigenvalue Problem. Clarendon Press, Oxford, 1965.

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J.H. WILKINSON - THE ALGEBRAIC EIGENVALUE PROBLEM	ATEI 520
.....	ATEI 530
CLARENDON PRESS, OXFORD, 1965.	ATEI 540

		ATEI 430
	SUBROUTINE ATEIG(M,A,RR,R1,IANA,IA)	ATEI 440
	CEPENSICA A(1),NR(1),RI(1),PRR(2),PRI(2),IAN(1)	ATEI 450
	INTEGER P,P1,C	ATEI 460
		ATEI 470
	E7=1,C6=8	ATEI 480
	E6=1,C6=6	ATEI 490
	E1C=1,C6=1C	ATEI 500
	CELTAC=C,5	ATEI 510
	PARIT=3C	ATEI 520
		ATEI 530
	INITIALIZATION	ATEI 540
		ATEI 550
	A=0	ATEI 560
	2C A1=A-1	ATEI 570
	IA=IA+1A	ATEI 580
	AA=IA+A	ATEI 590
	IF(IA) 3C,13CC,3C	ATEI 600
	3C AF=AA+1	ATEI 610
		ATEI 620
	ITERATION COUNTER	ATEI 630
		ATEI 640
	IT=C	ATEI 650
		ATEI 660
	RCCTS OF THE 2ND ORDER MAIN SUBMATRIX AT THE PREVIOUS	ATEI 670
	ITERATION	ATEI 680
		ATEI 690
	CC 4C I=1,2	ATEI 700
	PRR(1)=C,C	ATEI 710
	4C FFI(1)=C,C	ATEI 720
		ATEI 730
	LAST TWO SUBDIAGONAL ELEMENTS AT THE PREVIOUS ITERATION	ATEI 740
		ATEI 750
	FA=C,C	ATEI 760
	FA1=C,C	ATEI 770
		ATEI 780
	CRICIA SHIFT	ATEI 790
		ATEI 800
	R=C,C	ATEI 810
	S=C,C	ATEI 820
		ATEI 830
	RCCTS OF THE LOWER MAIN 2 BY 2 SUBMATRIX	ATEI 840
		ATEI 850
	A2=AA-1	ATEI 860
	IA1=IA-1A	ATEI 870
	AA1=IA1+A	ATEI 880
	IA1=IA+AA1	ATEI 890
	AA1=IA1+AA1	ATEI 900
	6C T=A(IA1)-A(IA)	ATEI 910
	L=T+0	ATEI 920
	V=4,C*A(IA1)+A(IA1)	ATEI 930
	IF(AES(V)-L+E7) 1CC,10C,65	ATEI 940
	65 T=V+V	ATEI 950
	IF(AES(T)-A*AA1(L,AES(V)))E61 E7,67,68	ATEI 960
	67 T=C,C	ATEI 970
	6E L=A(IA1)+A(IA1))/2,C	ATEI 980
	V=SQRT(AES(T))/2,C	ATEI 990
	IF(T)14C,7C,7C	ATEI 1000
	7C IF(L) 6C,75,75	ATEI 1010
	75 PR(1)=L+V	ATEI 1020
	RR(1)=L+V	ATEI 1030
	CC TC 13C	ATEI 1040
	6C RR(1)=L+V	ATEI 1050
	RR(1)=L+V	ATEI 1060
	CC TC 13C	ATEI 1070
	1CC IF(1)12C,11C,11C	ATEI 1080
	11C RR(1)=A(IA1)	ATEI 1090
	RR(1)=A(IA)	ATEI 1100
	CC TC 13C	ATEI 1110
	12C PR(1)=A(IA)	ATEI 1120
	RR(1)=A(IA1)	ATEI 1130
	13C R(1)=C,C	ATEI 1140
	R(1)=C,C	ATEI 1150
	CC TC 13C	ATEI 1160
	14C PR(1)=L	ATEI 1170
	RR(1)=L	ATEI 1180
	R(1)=V	ATEI 1190
	R(1)=V	ATEI 1200
	16C IF(A2)126C,126C,18C	ATEI 1210
		ATEI 1220
	TESTS OF CONVERGENCE	ATEI 1230
		ATEI 1240
	18C A1A2=AA1-1A	ATEI 1250
	MPCC=PR(1)+RR(1)+R(1)+R(1)	ATEI 1260
	EPS=E1C*SGT(MPCC)	ATEI 1270
	IF(AES(A1A2))-EPS)126C,126C,24C	ATEI 1280
	24C IF(AES(A1A1))-E1C+AES(A1A1)) 13CC,13CC,250	ATEI 1290
	25C IF(AES(A1A1)-A1A12)-AES(A1A12))E61 124C,124C,26C	ATEI 1300
	26C IF(AES(A1A1)-A1A11)-AES(A1A11))E61 124C,124C,30C	ATEI 1310
	30C IF(1)-PARIT) 32C,124C,124C	ATEI 1320
		ATEI 1330
	COMPLET THE SHIFT	ATEI 1340
		ATEI 1350
	32C J=1	ATEI 1360
	CC 36C I=1,2	ATEI 1370
	RR=1	ATEI 1380
	IF(AES(PRI(K)-PRR(1)+ABS(RI(K)-PRI(1))-CELTAC*ABS(PRI(K))	ATEI 1390
	1 -ABS(RI(K)))) 34C,36C,36C	ATEI 1400
	34C J=J+1	ATEI 1410
	36C CE1TAC=	ATEI 1420
	CC TC (44C,46C,46C,46C),J	ATEI 1430
	44C R=C,C	ATEI 1440
	S=C,C	ATEI 1450
	CC TC 5CC	ATEI 1460
	46C J=AA2-J	ATEI 1470
	FA=RR(J)+RR(J)	ATEI 1480
	S=RR(J)+RR(J)	ATEI 1490
	CC TC 5CC	ATEI 1500
	48C RR=RR(J)+RR(1)-RI(1)+RI(1)	ATEI 1510
	S=RR(J)+RR(1)	ATEI 1520
		ATEI 1530
		ATEI 1540
	SAVE THE LAST TWO SUBDIAGONAL TERMS AND THE RCCTS OF THE	ATEI 1550
	SUBMATRIX BEFORE ITERATION	ATEI 1560
		ATEI 1570
	50C FA=AA(1)	ATEI 1580
	FA1=A(IA2)	ATEI 1590
	CC 52C I=1,2	ATEI 1600
	RR=1	ATEI 1610
	PRR(1)=RR(1)	ATEI 1620
	52C PRI(1)=RI(K)	ATEI 1630
		ATEI 1640
	SEARCH FOR A PARTITION OF THE MATRIX, DEFINED BY P AND Q	ATEI 1650
		ATEI 1660
	FA=2	ATEI 1670
	IF (A-2)6CC,6CC,525	ATEI 1680
	525 IF(A1A2)	ATEI 1690
	CC 56C J=2,AA2	ATEI 1700
	IF(1)-PRI(1)-1	ATEI 1710
	IF(AES(A1PRI(1))-EPS) 6CC,6CC,53C	ATEI 1720

[illegible]

ATE11720
ATE11730
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ATE11780
ATE11790
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ATE11900
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ATE11970
ATE11980
ATE11990
ATE12000
ATE12010
ATE12020
ATE12030
ATE12040
ATE12050
ATE12060
ATE12070
ATE12080
ATE12090
ATE12100
ATE12110
ATE12120
ATE12130
ATE12140
ATE12150
ATE12160
ATE12170
ATE12180
ATE12190
ATE12200
ATE12210
ATE12220
ATE12230
ATE12240
ATE12250
ATE12260
ATE12270
ATE12280
ATE12290
ATE12300
ATE12310
ATE12320
ATE12330
ATE12340
ATE12350
ATE12360
ATE12370
ATE12380
ATE12390
ATE12400
ATE12410
ATE12420
ATE12430
ATE12440
ATE12450
ATE12460
ATE12470
ATE12480
ATE12490
ATE12500
ATE12510
ATE12520
ATE12530
ATE12540
ATE12550
ATE12560
ATE12570
ATE12580
ATE12590
ATE12600
ATE12610
ATE12620
ATE12630
ATE12640
ATE12650
ATE12660
ATE12670
ATE12680
ATE12690
ATE12700
ATE12710
ATE12720
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ATE12890
ATE12900
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ATE12920
ATE12930
ATE12940
ATE12950
ATE12960
ATE12970
ATE12980
ATE12990

Subroutine HSBG

This subroutine reduces an n by n real matrix A by a similarity transformation to upper almost-triangular (Hessenberg) form. Each row is reduced in turn, starting from the last one, by applying a suitable right elimination matrix, and similarity is achieved by also applying the left inverse transformation. Thus the eigenvalues of A are preserved.

1. Mathematical background

Let $A^{(p)}$ denote the matrix obtained from $A^{(0)} = A$ after reducing rows $n, n-1, \dots, n-p+1$. The similarity which transforms $A^{(p)}$ to $A^{(p+1)}$ is as follows:

a. First we determine a pivot element $a_{n-p, k}^{(p)}$ whose column subscript k is such that

$$\left| a_{n-p, k}^{(p)} \right| = \max_i \left| a_{n-p, i}^{(p)} \right|, i = 1, 2, \dots, n-p-1$$

b. If the pivot element $a_{n-p, k}^{(p)} = 0$, no transformations are necessary; that is, the $(p+1)$ th similarity is the identity transformation. Otherwise, if it is necessary ($k \neq n-p-1$), we interchange the k th and $(n-p-1)$ th columns so that the pivot element is in the subdiagonal position. The same interchange is applied to the rows of $A^{(p)}$. The resulting matrix is similar to $A^{(p)}$, and, to ease the notation, we denote it by $A^{(p)}$.

c. Define multipliers:

$$b_j^{(p)} = a_{n-p,j}^{(p)} / a_{n-p,n-p-1}^{(p)} \quad j=1, 2, \dots, n-p-2$$

Then the $(p+1)$ th similarity is given by the following.

Right elimination:

$$\bar{a}_{ij}^{(p)} = a_{ij}^{(p)} - b_j^{(p)} a_{i, n-p-1}^{(p)}$$

$$\begin{aligned} i &= 1, 2, \dots, n-p \\ j &= 1, 2, \dots, n-p-2 \end{aligned}$$

$$\bar{a}_{ij}(p) = a_{ij}(p) \quad \text{other indices}$$

Left inverse transformation:

$$a_{n-p-1,j}^{(p+1)} = \bar{a}_{n-p-1,j}^{(p)} + \sum_{i=1}^{n-p-2} b_i^{(p)} \bar{a}_{ij}^{(p)}$$

$$a_{ij}^{(p+1)} = \bar{a}_{ij}^{(p)} \quad i \neq n-p-1$$

Finally, $A^{(n-2)}$ will have the upper almost-triangular form.

For reference see J. H. Wilkinson, The Algebraic Eigenvalue Problem. Clarendon Press, Oxford, 1965.

```

K=K+IA
LK=L+L1
S=A(LK)
LJ=L-IA
DO 280 J=1,L2
JK=K+J
LJ=L+JA
280 S=S+A(LJ)*A(LK)*1.000
300 A(LK)=S
C
C      SET THE LOWER PART OF THE MATRIX TO ZERO
C
DO 310 I=L,LIA,IA
310 A(I)=0.0
320 L=L1
GO TO 20
360 RETURN
END

```

```

H5RG1070
H5RG1080
H5RG1090
H5RG1100
H5RG1110
H5RG1120
H5RG1130
H5RG1140
H5RG1150
H5RG1160
H5RG1170
H5RG1180
H5RG1190
H5RG1200
H5RG1210
H5RG1220
H5RG1230
H5RG1240

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C .....
C      SUBROUTINE HSBG
C      PURPOSE
C      TO REDUCE A REAL MATRIX INTO UPPER ALMOST TRIANGULAR FORM
C      USAGE
C      CALL HSBG(N,A,IA)
C      DESCRIPTION OF THE PARAMETERS
C      N      ORDER OF THE MATRIX
C      A      THE INPUT MATRIX, N BY N
C      IA     SIZE OF THE FIRST DIMENSION ASSIGNED TO THE ARRAY
C      A IN THE CALLING PROGRAM WHEN THE MATRIX IS IN
C      DOUBLE SUBSCRIPTED DATA STORAGE MODE, IA=N WHEN
C      THE MATRIX IS IN SSP VECTOR STORAGE MODE.
C      REMARKS
C      THE HESSENBERG FORM REPLACES THE ORIGINAL MATRIX IN THE
C      ARRAY A.
C      SUBROUTINES AND FUNCTION SUBPROGRAMS REQUIRED
C      NONE
C      METHOD
C      SIMILARITY TRANSFORMATIONS USING ELEMENTARY ELIMINATION
C      MATRICES, WITH PARTIAL PIVOTING.
C      REFERENCES
C      J.H. WILKINSON - THE ALGEBRAIC EIGENVALUE PROBLEM -
C      CLARENDON PRESS, OXFORD, 1965.
C .....
C      SUBROUTINE HSBG(N,A,IA)
C      DIMENSION A(1)
C      DOUBLE PRECISION S
C      L=N
C      NIA=L-IA
C      LIA=NIA-IA
C      L IS THE ROW INDEX OF THE ELIMINATION
C
C      20 IF(L-3) 360,40,40
C      40 LIA=LIA-IA
C      LI=L-1
C      L2=L1-1
C
C      SEARCH FOR THE PIVOTAL ELEMENT IN THE LTH ROW
C
C      ISUB=LIA+L
C      IPIV=ISUB-IA
C      PIV=ABS(A(IPIV))
C      IF(L-3) 90,40,50
C      50 M=IPIV-IA
C      DO 90 I=L,M,IA
C      1=ABS(A(I))
C      IF(1>PIV) 80,80,60
C      60 IPIV=I
C      80 CONTINUE
C      90 IF(PIV) 100,120,100
C      100 IF(PIV=ABS(A(ISUB))) 180,180,120
C
C      INTERCHANGE THE COLUMNS
C
C      120 M=IPIV-L
C      DO 140 I=L,L
C      J=M+I
C      T=A(J)
C      K=LIA+I
C      A(J)=A(K)
C      140 A(K)=T
C
C      INTERCHANGE THE ROWS
C
C      M=L2-M/IA
C      DO 160 I=L1,NIA,IA
C      T=A(I)
C      J=L-M
C      A(I)=A(J)
C      160 A(J)=T
C
C      TERMS OF THE ELEMENTARY TRANSFORMATION
C
C      180 DO 200 I=L,LIA,IA
C      200 A(I)=A(I)/A(ISUB)
C
C      RIGHT TRANSFORMATION
C
C      J=-IA
C      DO 240 I=1,L2
C      J=J+IA
C      LJ=L+J
C      DO 220 K=1,L1
C      KJ=K+J
C      KL=K-LIA
C      220 A(KJ)=A(KJ)-A(LJ)*A(KL)
C      240 CONTINUE
C
C      LEFT TRANSFORMATION
C
C      K=-IA
C      DO 300 I=1,N

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```

H5BG 10
H5BG 20
H5BG 30
H5BG 40
H5BG 50
H5BG 60
H5BG 70
H5BG 80
H5BG 90
H5BG 100
H5BG 110
H5BG 120
H5BG 130
H5BG 140
H5BG 150
H5BG 160
H5BG 170
H5BG 180
H5BG 190
H5BG 200
H5BG 210
H5BG 220
H5BG 230
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H5BG 320
H5BG 330
H5BG 340
H5BG 350
H5BG 360
H5BG 370
H5BG 380
H5BG 390
H5BG 400
H5BG 410
H5BG 420
H5BG 430
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H5BG 450
H5BG 460
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H5BG 580
H5BG 590
H5BG 600
H5BG 610
H5BG 620
H5BG 630
H5BG 640
H5BG 650
H5BG 660
H5BG 670
H5BG 680
H5BG 690
H5BG 700
H5BG 710
H5BG 720
H5BG 730
H5BG 740
H5BG 750
H5BG 760
H5BG 770
H5BG 780
H5BG 790
H5BG 800
H5BG 810
H5BG 820
H5BG 830
H5BG 840
H5BG 850
H5BG 860
H5BG 870
H5BG 880
H5BG 890
H5BG 900
H5BG 910
H5BG 920
H5BG 930
H5BG 940
H5BG 950
H5BG 960
H5BG 970
H5BG 980
H5BG 990
H5BG1000
H5BG1010
H5BG1020
H5BG1030
H5BG1040
H5BG1050
H5BG1060

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